

**PROBABILISTIC METHODS IN COMBINATORICS**  
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**PROBLEM SET**

A. INTRODUCTION AND LINEARITY OF EXPECTATIONS

A1. Verify the following asymptotic calculations used in Ramsey number lower bounds:

(a) For each  $k$ , the largest  $n$  satisfying  $\binom{n}{k}2^{1-\binom{k}{2}} < 1$  has  $n = \left(\frac{1}{e\sqrt{2}} + o(1)\right) k2^{k/2}$ .

(b) For each  $k$ , the maximum value of  $n - \binom{n}{k}2^{1-\binom{k}{2}}$  as  $n$  ranges over positive integers is  $\left(\frac{1}{e} + o(1)\right) k2^{k/2}$ .

(c) For each  $k$ , the largest  $n$  satisfying  $e \left(\binom{k}{2} \binom{n}{k-2} + 1\right) 2^{1-\binom{k}{2}} < 1$  satisfies  $n = \left(\frac{\sqrt{2}}{e} + o(1)\right) k2^{k/2}$ .

A2. Prove that, if there is a real  $p \in [0, 1]$  such that

$$\binom{n}{k} p^{\binom{k}{2}} + \binom{n}{t} (1-p)^{\binom{t}{2}} < 1$$

then the Ramsey number  $R(k, t)$  satisfies  $R(k, t) > n$ . Using this show that

$$R(4, t) \geq c \left(\frac{t}{\log t}\right)^{3/2}$$

for some constant  $c > 0$ .

ps1

A3. Let  $G$  be a graph with  $n$  vertices and  $m$  edges. Prove that  $K_n$  can be written as a union of  $O(n^2(\log n)/m)$  isomorphic copies of  $G$  (not necessarily edge-disjoint).

A4. Prove that there is an absolute constant  $C > 0$  so that for every  $n \times n$  matrix with distinct real entries, one can permute its rows so that no column in the permuted matrix contains an increasing subsequence of length at least  $C\sqrt{n}$ . (A subsequence does not need to be selected from consecutive terms. For example,  $(1, 2, 3)$  is an increasing subsequence of  $(1, 5, 2, 4, 3)$ .)

A5. *Generalization of Sperner's theorem.* Let  $\mathcal{F}$  be a collection of subset of  $[n]$  that does not contain  $k + 1$  elements forming a chain:  $A_1 \subsetneq \dots \subsetneq A_{k+1}$ . Prove that  $\mathcal{F}$  is no larger than taking the union of the  $k$  levels of the Boolean lattice closest to the middle layer.

A6. Let  $G$  be a graph on  $n \geq 10$  vertices. Suppose that adding any new edge to  $G$  would create a new clique on 10 vertices. Prove that  $G$  has at least  $8n - 36$  edges.

Hint in white:

A7. Let  $k \geq 4$  and  $H$  a  $k$ -uniform hypergraph with at most  $4^{k-1}/3^k$  edges. Prove that there is a coloring of the vertices of  $H$  by four colors so that in every edge all four colors are represented.

ps1

A8. Given a set  $\mathcal{F}$  of subsets of  $[n]$  and  $A \subseteq [n]$ , write  $\mathcal{F}|_A := \{S \cap A : S \in \mathcal{F}\}$  (its *projection* onto  $A$ ). Prove that for every  $n$  and  $k$ , there exists a set  $\mathcal{F}$  of subsets of  $[n]$  with  $|\mathcal{F}| = O(k2^k \log n)$  such that for every  $k$ -element subset  $A$  of  $[n]$ ,  $\mathcal{F}|_A$  contains all  $2^k$  subsets of  $A$ .

ps1

A9. Let  $A_1, \dots, A_m$  be  $r$ -element sets and  $B_1, \dots, B_m$  be  $s$ -element sets. Suppose  $A_i \cap B_i = \emptyset$  for each  $i$ , and for each  $i \neq j$ , either  $A_i \cap B_j \neq \emptyset$  or  $A_j \cap B_i \neq \emptyset$ . Prove that  $m \leq (r+s)^{r+s}/(r^r s^s)$ .

- ps1★ A10. Show that in every non-2-colorable  $n$ -uniform hypergraph, one can find at least  $\frac{n}{2} \binom{2n-1}{n}$  unordered pairs of edges with each pair intersecting in exactly one vertex.
- A11. Let  $A$  be a measurable subset of the unit sphere in  $\mathbb{R}^3$  (centered at the origin) containing no pair of orthogonal points.
- ps1 (a) Prove that  $A$  occupies at most  $1/3$  of the sphere in terms of surface area.
- ps1★ (b) Prove an upper bound smaller than  $1/3$  (give your best bound).
- ps1★ A12. Prove that every set of 10 points in the plane can be covered by a union of disjoint unit disks.
- A13. Let  $\mathbf{r} = (r_1, \dots, r_k)$  be a vector of nonzero integers whose sum is nonzero. Prove that there exists a real  $c > 0$  (depending on  $\mathbf{r}$  only) such that the following holds: for every finite set  $A$  of nonzero reals, there exists a subset  $B \subseteq A$  with  $|B| \geq c|A|$  such that there do not exist  $b_1, \dots, b_k \in B$  with  $r_1 b_1 + \dots + r_k b_k = 0$ .
- ps1 A14. Prove that every set  $A$  of  $n$  nonzero integers contains two disjoint subsets  $B_1$  and  $B_2$ , such that both  $B_1$  and  $B_2$  are sum-free, and  $|B_1| + |B_2| > 2n/3$ .
- ps1 A15. Let  $G$  be an  $n$ -vertex graph with  $pn^2$  edges, with  $n \geq 10$  and  $p \geq 10/n$ . Prove that  $G$  contains a pair of vertex-disjoint and isomorphic subgraphs (not necessarily induced) each with at least  $cp^2n^2$  edges, where  $c > 0$  is a constant.
- ps1★ A16. Prove that for every positive integer  $r$ , there exists an integer  $K$  such that the following holds. Let  $S$  be a set of  $rk$  points evenly spaced on a circle. If we partition  $S = S_1 \cup \dots \cup S_r$  so that  $|S_i| = k$  for each  $i$ , then, provided  $k \geq K$ , there exist  $r$  congruent triangles where the vertices of the  $i$ -th triangle lie in  $S_i$ , for each  $1 \leq i \leq r$ .
- ps1★ A17. Prove that  $[n]^d$  cannot be partitioned into fewer than  $2^d$  sets each of the form  $A_1 \times \dots \times A_d$  where  $A_i \subsetneq [n]$ .

B. ALTERATION METHOD

- B1. Using the alteration method, prove the Ramsey number bound

$$R(4, k) \geq c(k/\log k)^2$$

for some constant  $c > 0$ .

- B2. Prove that every 3-uniform hypergraph with  $n$  vertices and  $m \geq n$  edges contains an independent set (i.e., a set of vertices containing no edges) of size at least  $cn^{3/2}/\sqrt{m}$ , where  $c > 0$  is a constant.
- B3. Prove that every  $k$ -uniform hypergraph with  $n$  vertices and  $m$  edges has a transversal (i.e., a set of vertices intersecting every edge) of size at most  $n(\log k)/k + m/k$ .
- ps2 B4. *Zarankiewicz problem*. Prove that for every positive integers  $n \geq k \geq 2$ , there exists an  $n \times n$  matrix with  $\{0, 1\}$  entries, with at least  $\frac{1}{2}n^{2-2/(k+1)}$  1's, such that there is no  $k \times k$  submatrix consisting of all 1's.
- ps2 B5. Fix  $k$ . Prove that there exists a constant  $c_k > 1$  so that for every sufficiently large  $n > n_0(k)$ , there exists a collection  $\mathcal{F}$  of at least  $c_k^n$  subsets of  $[n]$  such that for every  $k$  distinct  $F_1, \dots, F_k \in \mathcal{F}$ , all  $2^k$  intersections  $\bigcap_{i=1}^k G_i$  are nonempty, where each  $G_i$  is either  $F_i$  or  $[n] \setminus F_i$ .

B6. *Acute sets in  $\mathbb{R}^n$ .* Prove that, for some constant  $c > 0$ , for every  $n$ , there exists a family of at least  $c(2/\sqrt{3})^n$  subsets of  $[n]$  containing no three distinct members  $A, B, C$  satisfying  $A \cap B \subseteq C \subseteq A \cup B$ .

Deduce that there exists a set of at least  $c(2/\sqrt{3})^n$  points in  $\mathbb{R}^n$  so that all angles determined by three points from the set are acute.

*Remark.* The current best lower and upper bounds for the maximum size of an “acute set” in  $\mathbb{R}^n$  (i.e., spanning only acute angles) are  $2^{n-1} + 1$  and  $2^n - 1$  respectively.

ps2★

B7. *Covering complements of sparse graphs by cliques*

(a) Prove that every graph with  $n$  vertices and minimum degree  $n - d$  can be written as a union of  $O(d^2 \log n)$  cliques.

(b) Prove that every bipartite graph with  $n$  vertices on each side of the vertex bipartition and minimum degree  $n - d$  can be written as a union of  $O(d \log n)$  complete bipartite graphs (assume  $d \geq 1$ ).

ps2★

B8. Let  $G = (V, E)$  be a graph with  $n$  vertices and minimum degree  $\delta \geq 2$ . Prove that there exists  $A \subseteq V$  with  $|A| = O(n(\log \delta)/\delta)$  so that every vertex in  $V \setminus A$  contains at least one neighbor in  $A$  and at least one neighbor not in  $A$ .

ps2★

B9. Prove that every graph  $G$  without isolated vertices has an induced subgraph  $H$  on at least  $\alpha(G)/2$  vertices such that all vertices of  $H$  have odd degree. Here  $\alpha(G)$  is the size of the largest independent set in  $G$ .

### C. SECOND MOMENT METHOD

ps2

C1. *Threshold for  $k$ -APs.* Let  $[n]_p$  denote the random subset of  $\{1, \dots, n\}$  where every element is included with probability  $p$  independently. For each fixed integer  $k \geq 3$ , determine the threshold for  $[n]_p$  to contain a  $k$ -term arithmetic progression.

C2. Show that, for each fixed positive integer  $k$ , there is a sequence  $p_n$  such that

$$\mathbb{P}(G(n, p_n) \text{ has a connected component with exactly } k \text{ vertices}) \rightarrow 1 \quad \text{as } n \rightarrow \infty.$$

Hint in white:

ps2

C3. *Poisson limit.* Let  $X$  be the number of triangles in  $G(n, c/n)$  for some fixed  $c > 0$ .

(a) For every nonnegative integer  $k$ , determine the limit of  $\mathbb{E}\binom{X}{k}$  as  $n \rightarrow \infty$ .

(b) Let  $Y \sim \text{Binomial}(n, \lambda/n)$  for some fixed  $\lambda > 0$ . For every nonnegative integer  $k$ , determine the limit of  $\mathbb{E}\binom{Y}{k}$  as  $n \rightarrow \infty$ , and show that it agrees with the limit in (a) for some  $\lambda = \lambda(c)$ .

We know that  $Y$  converges to the Poisson distribution with mean  $\lambda$ . Also, the Poisson distribution is determined by its moments.

(c) Compute, for fixed every nonnegative integer  $t$ , the limit of  $\mathbb{P}(X = t)$  as  $n \rightarrow \infty$ .

(In particular, this gives the limit probability that  $G(n, c/n)$  contains a triangle, i.e.,  $\lim_{n \rightarrow \infty} \mathbb{P}(X > 0)$ . This limit increases from 0 to 1 continuously when  $c$  ranges from 0 to  $+\infty$ , thereby showing that the property of containing a triangle has a coarse threshold.)

ps2

C4. *Central limit theorem for triangle counts.* Find a real (non-random) sequence  $a_n$  so that, letting  $X$  be the number of triangles and  $Y$  be the number of edges in the random graph

$G(n, 1/2)$ , one has

$$\text{Var}(X - a_n Y) = o(\text{Var } X).$$

Deduce that  $X$  is asymptotically normal, that is,  $(X - \mathbb{E}X)/\sqrt{\text{Var } X}$  converges to the normal distribution.

(You can solve for the minimizing  $a_n$  directly similar to ordinary least squares linear regression, or first write the edge indicator variables as  $X_{ij} = (1 + Y_{ij})/2$  and then expand. The latter approach likely yields a cleaner computation.)

C5. *Isolated vertices.* Let  $p_n = (\log n + c_n)/n$ .

(a) Show that, as  $n \rightarrow \infty$ ,

$$\mathbb{P}(G(n, p_n) \text{ has no isolated vertices}) \rightarrow \begin{cases} 0 & \text{if } c_n \rightarrow -\infty, \\ 1 & \text{if } c_n \rightarrow \infty. \end{cases}$$

(b) Suppose  $c_n \rightarrow c \in \mathbb{R}$ , compute, with proof, the limit of LHS above as  $n \rightarrow \infty$ , by following the approach in **C3**.

ps2\*

C6. Is the threshold for the bipartiteness of a random graph coarse or sharp?

(You are not allowed to use any theorems that we did not prove in class/notes.)

ps2

C7. *Triangle packing.* Prove that, with probability approaching 1 as  $n \rightarrow \infty$ ,  $G(n, n^{-1/2})$  has at least  $cn^{3/2}$  edge-disjoint triangles, where  $c > 0$  is some constant.

Hint in white:

ps3

C8. *Nearly perfect triangle factor.* Prove that, with probability approaching 1 as  $n \rightarrow \infty$ ,

(a)  $G(n, n^{-2/3})$  has at least  $n/100$  vertex-disjoint triangles.

(b) *Simple nibble.*  $G(n, Cn^{-2/3})$  has at least  $0.33n$  vertex-disjoint triangles, for some constant  $C$ .

Hint in white:

C9. *Permuted correlation.* Recall that the *correlation* of two non-constant random variables  $X$  and  $Y$  is defined to be  $\text{corr}(X, Y) := \text{Cov}[X, Y]/\sqrt{(\text{Var } X)(\text{Var } Y)}$ .

Let  $f, g \in [n] \rightarrow \mathbb{R}$  be two non-constant functions. Prove that there exist permutations  $\pi$  and  $\tau$  of  $[n]$  such that, with  $Z$  being a uniform random element of  $[n]$ ,

$$\text{corr}(f(\pi(Z)), g(Z)) - \text{corr}(f(\tau(Z)), g(Z)) \geq \frac{2}{\sqrt{n-1}}.$$

Furthermore, show that equality can be achieved for even  $n$ .

Hint in white:

ps3

C10. Let  $v_1 = (x_1, y_1), \dots, v_n = (x_n, y_n) \in \mathbb{Z}^2$  with  $|x_i|, |y_i| \leq 2^{n/2}/(100\sqrt{n})$  for all  $i \in [n]$ . Show that there are two disjoint sets  $I, J \subseteq [n]$ , not both empty, such that  $\sum_{i \in I} v_i = \sum_{j \in J} v_j$ .

ps3\*

C11. Prove that there is an absolute constant  $C > 0$  so that the following holds. For every prime  $p$  and every  $A \subseteq \mathbb{Z}/p\mathbb{Z}$  with  $|A| = k$ , there exists an integer  $x$  so that  $\{xa : a \in A\}$  intersects every interval of length at least  $Cp/\sqrt{k}$  in  $\mathbb{Z}/p\mathbb{Z}$ .

ps3\*

C12. Prove that there is a constant  $c > 0$  so that every hyperplane containing the origin in  $\mathbb{R}^n$  intersects at least  $c$ -fraction of the  $2^n$  closed unit balls centered at  $\{-1, 1\}^n$ .

## D. CHERNOFF BOUND

D1. Prove that with probability  $1 - o(1)$  as  $n \rightarrow \infty$ , every bipartite subgraph of  $G(n, 1/2)$  has at most  $n^2/8 + 10n^{3/2}$  edges.

ps3

D2. *Unbalancing lights.* Prove that there is a constant  $C$  so that for every positive integer  $n$ , one can find an  $n \times n$  matrix  $A$  with  $\{-1, 1\}$  entries, so that for all vectors  $x, y \in \{-1, 1\}^n$ ,  $|y^\top Ax| \leq Cn^{3/2}$ .

ps3

D3. Prove that there exists a constant  $c > 1$  such that for every  $n$ , there are at least  $c^n$  points in  $\mathbb{R}^n$  so that every triple of points form a triangle whose angles are all less than  $61^\circ$ .

ps3

D4. *Planted clique.* Give a deterministic polynomial-time algorithm for the following task so that it succeeds over the random input with probability approaching 1 as  $n \rightarrow \infty$ .

Input: some unlabeled  $n$ -vertex  $G$  created as the union of  $G(n, 1/2)$  and a clique on  $t = \lfloor 100\sqrt{n \log n} \rfloor$  vertices.

Output: a clique in  $G$  of size  $t$ .

D5. *Weighing coins.* You are given  $n$  coins, each with one of two known weights, but otherwise indistinguishable. You can use a scale that outputs the combined weight of any subset of the coins. You must decide in advance which subsets  $S_1, \dots, S_k \subseteq [n]$  of the coins to weigh. We wish to determine the minimum number of weighings needed to identify the weight of every coin. (Below,  $X$  and  $Y$  represent two possibilities for which coins are of the first weight.)

ps3\*

(a) Prove that if  $k \leq 1.99n/\log_2 n$  and  $n$  is sufficiently large, then for every  $S_1, \dots, S_k \subseteq [n]$ , there are two distinct subsets  $X, Y \subseteq [n]$  such that  $|X \cap S_i| = |Y \cap S_i|$  for all  $i \in [k]$ .

(There is a neat solution to part (a) using information theory, though here you are explicitly asked to solve it using the Chernoff bound.)

ps3\*

(b) Show that there is some constant  $C$  such that (a) is false if 1.99 is replaced by  $C$ . (What is the best  $C$  you can get?)

## E. LOVÁSZ LOCAL LEMMA

ps3

E1. Show that it is possible to color the edges of  $K_n$  with at most  $3\sqrt{n}$  colors so that there are no monochromatic triangles.

E2. Prove that it is possible to color the vertices of every  $k$ -uniform  $k$ -regular hypergraph using at most  $k/\log k$  colors so that every color appears at most  $O(\log k)$  times on each edge.

ps3\*

E3. *Hitting thin rectangles.* Prove that there is a constant  $C > 0$  so that for every sufficiently small  $\epsilon > 0$ , one can choose exactly one point inside each grid square  $[n, n+1) \times [m, m+1) \subset \mathbb{R}^2$ ,  $m, n \in \mathbb{Z}$ , so that every rectangle of dimensions  $\epsilon$  by  $(C/\epsilon) \log(1/\epsilon)$  in the plane (not necessarily axis-aligned) contains at least one chosen point.

ps4

E4. *List coloring.* Prove that there is some constant  $c > 0$  so that given a graph and a set of  $k$  acceptable colors for each vertex such that every color is acceptable for at most  $ck$  neighbors of each vertex, there is always a proper coloring where every vertex is assigned one of its acceptable colors.

E5. Prove that, for every  $\epsilon > 0$ , there exist  $\ell_0$  and some  $(a_1, a_2, \dots) \in \{0, 1\}^{\mathbb{N}}$  such that for every  $\ell > \ell_0$  and every  $i > 1$ , the vectors  $(a_i, a_{i+1}, \dots, a_{i+\ell-1})$  and  $(a_{i+\ell}, a_{i+\ell+1}, \dots, a_{i+2\ell-1})$  differ in at least  $(\frac{1}{2} - \epsilon)\ell$  coordinates.

ps4 E6. *Avoiding periodically colored paths.* Prove that for every  $\Delta$ , there exists  $k$  so that every graph with maximum degree at most  $\Delta$  has a vertex-coloring using  $k$  colors so that there is no path of the form  $v_1v_2 \dots v_{2\ell}$  (for any positive integer  $\ell$ ) where  $v_i$  has the same color as  $v_{i+\ell}$  for each  $i \in [\ell]$ . (Note that vertices on a path must be distinct.)

ps4 E7. Prove that every graph with maximum degree  $\Delta$  can be properly edge-colored using  $O(\Delta)$  colors so that every cycle contains at least three colors.

(An edge-coloring is *proper* if it never assigns the same color to two edges sharing a vertex.)

ps4\* E8. Prove that for every  $\Delta$ , there exists  $g$  so that every bipartite graph with maximum degree  $\Delta$  and girth at least  $g$  can be properly edge-colored using  $\Delta + 1$  colors so that every cycle contains at least three colors.

ps4\* E9. Prove that for every positive integer  $r$ , there exists  $C_r$  so that every graph with maximum degree  $\Delta$  has a *proper* vertex coloring using at most  $C_r\Delta^{1+1/r}$  colors so that every vertex has at most  $r$  neighbors of each color.

E10. *Vertex-disjoint cycles in digraphs.* (Recall that a directed graph is  $k$ -regular if all vertices have in-degree and out-degree both equal to  $k$ . Also, cycles cannot repeat vertices.)

ps4 (a) Prove that every  $k$ -regular directed graph has at least  $ck/\log k$  vertex-disjoint directed cycles, where  $c > 0$  is some constant.

ps4\* (b) Prove that every  $k$ -regular directed graph has at least  $ck$  vertex-disjoint directed cycles, where  $c > 0$  is some constant.

Hint in white:

E11. (a) *Generalization of Cayley's formula.* Using Prüfer codes, prove the identity

$$x_1x_2 \cdots x_n(x_1 + \cdots + x_n)^{n-2} = \sum_T x_1^{d_T(1)} x_2^{d_T(2)} \cdots x_n^{d_T(n)}$$

where the sum is over all trees  $T$  on  $n$  vertices labeled by  $[n]$  and  $d_T(i)$  is the degree of vertex  $i$  in  $T$ .

(b) Let  $F$  be a forest with vertex set  $[n]$ , with components having  $f_1, \dots, f_s$  vertices so that  $f_1 + \cdots + f_s = n$ . Prove that the number of trees on the vertex set  $[n]$  that contain  $F$  is exactly  $n^{n-2} \prod_{i=1}^s (f_i/n^{f_i-1})$ .

(c) *Independence property for uniform spanning tree of  $K_n$ .* Show that if  $H_1$  and  $H_2$  are vertex-disjoint subgraphs of  $K_n$ , then for a uniformly random spanning tree  $T$  of  $K_n$ , the events  $H_1 \subseteq T$  and  $H_2 \subseteq T$  are independent.

ps4\* (d) *Packing rainbow spanning trees.* Prove that there is a constant  $c > 0$  so that for every edge-coloring of  $K_n$  where each color appears at most  $cn$  times, there exist at least  $cn$  edge-disjoint spanning trees, where each spanning tree has all its edges colored differently.

(In your submission, you may assume previous parts without proof.)

*The next two problems use the lopsided local lemma.*

ps4 E12. *Packing two copies of a graph.* Prove that there is a constant  $c > 0$  so that if  $H$  is an  $n$ -vertex  $m$ -edge graph with maximum degree at most  $cn^2/m$ , then one can find two edge-disjoint copies of  $H$  in the complete graph  $K_n$ .

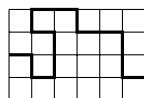
- ps4\* E13. *Packing Latin transversals.* Prove that there is a constant  $c > 0$  so that every  $n \times n$  matrix where no entry appears more than  $cn$  times contains  $cn$  disjoint Latin transversals.

F. CORRELATION INEQUALITIES

- F1. Let  $G = (V, E)$  be a graph. Color every edge with red or blue independently and uniformly at random. Let  $E_0$  be the set of red edges and  $E_1$  the set of blue edges. Let  $G_i = (V, E_i)$  for each  $i = 0, 1$ . Prove that

$$\mathbb{P}(G_0 \text{ and } G_1 \text{ are both connected}) \leq \mathbb{P}(G_0 \text{ is connected})^2.$$

- F2. A set family  $\mathcal{F}$  is *intersecting* if  $A \cap B \neq \emptyset$  for all  $A, B \in \mathcal{F}$ . Let  $\mathcal{F}_1, \dots, \mathcal{F}_k$  each be a collection of subsets of  $[n]$  and suppose that each  $\mathcal{F}_i$  is intersecting. Prove that  $\left| \bigcup_{i=1}^k \mathcal{F}_i \right| \leq 2^n - 2^{n-k}$ .
- F3. *Percolation.* Let  $G_{m,n}$  be the grid graph on vertex set  $[m] \times [n]$  ( $m$  vertices wide and  $n$  vertices tall). A *horizontal crossing* is a path that connects some left-most vertex to some right-most vertex. See below for an example of a horizontal crossing in  $G_{7,5}$ .



Let  $H_{m,n}$  denote the random subgraph of  $G_{m,n}$  obtained by keeping every edge with probability  $1/2$  independently.

Let  $\text{RSW}(k)$  denote the following statement: there exists a constant  $c_k > 0$  such that for all positive integers  $n$ ,  $\mathbb{P}(H_{kn,n} \text{ has a horizontal crossing}) \geq c_k$ .

ps5

- (a) Prove  $\text{RSW}(1)$ .

ps5

- (b) Prove that  $\text{RSW}(2)$  implies  $\text{RSW}(100)$ .

- (c) (Challenging) Prove  $\text{RSW}(2)$ .

- F4. Let  $A$  and  $B$  be two *independent* increasing events of independent random variables. Prove that there are two *disjoint* subsets  $S$  and  $T$  of these random variables so that  $A$  depends only on  $S$  and  $B$  depends only on  $T$ .

- F5. Let  $U_1$  and  $U_2$  be increasing events and  $D$  a decreasing event of independent Boolean random variables. Suppose  $U_1$  and  $U_2$  are independent. Prove that  $\mathbb{P}(U_1|U_2 \cap D) \leq \mathbb{P}(U_1|U_2)$ .

ps5

- F6. *Coupon collector.* Let  $s_1, \dots, s_m$  be independent random elements in  $[n]$  (not necessarily uniform or identically distributed; chosen with replacement) and  $S = \{s_1, \dots, s_m\}$ . Let  $I$  and  $J$  be disjoint subsets of  $[n]$ . Prove that  $\mathbb{P}(I \cup J \subseteq S) \leq \mathbb{P}(I \subseteq S)\mathbb{P}(J \subseteq S)$ .

ps5\*

- F7. Prove that there exist  $c < 1$  and  $\epsilon > 0$  such that if  $A_1, \dots, A_k$  are increasing events of independent Boolean random variables with  $\mathbb{P}(A_i) \leq \epsilon$  for all  $i$ , then, letting  $X$  denote the number of events  $A_i$  that occur, one has  $\mathbb{P}(X = 1) \leq c$ . (Give your smallest  $c$ . It is conjectured that any  $c > 1/e$  works.)

ps5\*

- F8. *Disjoint containment.* Let  $\mathcal{S}$  and  $\mathcal{T}$  each be a collection of subsets of  $[n]$ . Let  $R \subseteq [n]$  be a random subset where each element is included independently (not necessarily with the same probability). Let  $A$  be the event that  $S \subseteq R$  for some  $S \in \mathcal{S}$ . Let  $B$  be the event that  $T \subseteq R$  for some  $T \in \mathcal{T}$ . Let  $C$  denote the event there exist *disjoint*  $S, T \subseteq R$  with  $S \in \mathcal{S}$  and  $T \in \mathcal{T}$ . Prove that  $\mathbb{P}(C) \leq \mathbb{P}(A)\mathbb{P}(B)$ .

G. JANSON INEQUALITIES

- ps5 G1. *3-AP-free probability.* Determine, for all  $0 < p \leq 0.99$  ( $p$  is allowed to depend on  $n$ ), the probability that  $[n]_p$  does not contain a 3-term arithmetic progression, up to a constant factor in the exponent. (The form of the answer should be similar to the conclusion in class about the probability that  $G(n, p)$  is triangle-free. See C1 for notation.)
- G2. Prove that with probability  $1 - o(1)$ , the size of the largest subset of vertices of  $G(n, 1/2)$  inducing a triangle-free subgraph is  $\Theta(\log n)$ .
- G3. *Nearly perfect triangle factor, again.* Using Janson inequalities this time, give another solution to Problem C8 in the following generality.

- ps5 (a) Prove that for every  $\epsilon > 0$ , there exists  $C_\epsilon > 0$  such that with probability  $1 - o(1)$ ,  $G(n, C_\epsilon n^{-2/3})$  contains at least  $(1/3 - \epsilon)n$  vertex-disjoint triangles.
- (b) (Optional) Compare the dependence of the optimal  $C_\epsilon$  on  $\epsilon$  you obtain using the method in Problem C8 versus this problem (don't worry about leading constant factors).

- ps5★ G4. *Threshold for extensions.* Show that for every constant  $C > 16/5$ , if  $n^2 p^5 > C \log n$ , then with probability  $1 - o(1)$ , every edge of  $G(n, p)$  is contained in a  $K_4$ .

Be careful, this event is not increasing, and so it is insufficient to just prove the result for one specific  $p$ .

- G5. *Lower tails of small subgraph counts.* Fix graph  $H$  and  $\delta \in (0, 1]$ . Let  $X_H$  denote the number of copies of  $H$  in  $G(n, p)$ . Prove that for all  $n$  and  $0 < p < 0.99$ ,

$$\mathbb{P}(X_H \leq (1 - \delta)\mathbb{E}X_H) = e^{-\Theta_{H,\delta}(\Phi_H)} \quad \text{where } \Phi_H := \min_{H' \subseteq H: e(H') > 0} n^{v(H')} p^{e(H')}.$$

Here the hidden constants in  $\Theta_{H,\delta}$  may depend on  $H$  and  $\delta$  (but not on  $n$  and  $p$ ).

- ps5★ G6. *List chromatic number of a random graph.* Show that the list chromatic number of  $G(n, 1/2)$  is  $(1 + o(1))\frac{n}{2 \log_2 n}$  with probability  $1 - o(1)$ .

The *list-chromatic number* (also called *choosability*) of a graph  $G$  is defined to be the minimum  $k$  such that if every vertex of  $G$  is assigned a list of  $k$  acceptable colors, then there exists a proper coloring of  $G$  where every vertex is colored by one of its acceptable colors.

H. CONCENTRATION OF MEASURE

- ps5 H1. *Sub-Gaussian tails.* For each part, prove there is some constant  $c > 0$  so that, for all  $\lambda > 0$ ,

$$\mathbb{P}(|X - \mathbb{E}X| \geq \lambda \sqrt{\text{Var } X}) \leq 2e^{-c\lambda^2}.$$

- (a)  $X$  is the number of triangles in  $G(n, 1/2)$ .
- (b)  $X$  is the number of inversions of a uniform random permutation of  $[n]$  (an *inversion* of  $\sigma \in S_n$  is a pair  $(i, j)$  with  $i < j$  and  $\sigma(i) > \sigma(j)$ ).
- H2. Prove that for every  $\epsilon > 0$  there exists  $\delta > 0$  and  $n_0$  such that for all  $n \geq n_0$  and  $S_1, \dots, S_m \subset [2n]$  with  $m \leq 2^{\delta n}$  and  $|S_i| = n$  for all  $i \in [m]$ , there exists a function  $f: [2n] \rightarrow [n]$  so that  $(1 - e^{-1} - \epsilon)n \leq |f(S_i)| \leq (1 - e^{-1} + \epsilon)n$  for all  $i \in [m]$ .
- H3. *Simultaneous bisections.* Fix  $\Delta$ . Let  $G_1, \dots, G_m$  with  $m = 2^{o(n)}$  be connected graphs of maximum degree at most  $\Delta$  on the same vertex set  $V$  with  $|V| = n$ . Prove that there exists a partition  $V = A \cup B$  so that every  $G_i$  has  $(1 + o(1))e(G_i)/2$  edges between  $A$  and  $B$ .



- ps5\* H4. Prove that there is some constant  $c > 0$  so that for every graph  $G$  with chromatic number  $k$ , letting  $S$  be a uniform random subset of  $V$  and  $G[S]$  the subgraph induced by  $S$ , one has, for every  $t \geq 0$ ,

$$\mathbb{P}(\chi(G[S]) \leq k/2 - t) \leq e^{-ct^2/k}.$$

- ps5\* H5. Prove that there is some constant  $c > 0$  so that, with probability  $1 - o(1)$ ,  $G(n, 1/2)$  has a bipartite subgraph with at least  $n^2/8 + cn^{3/2}$  edges.

- H6. Let  $k \leq n/2$  be positive integers and  $G$  an  $n$ -vertex graph with average degree at most  $n/k$ . Prove that a uniform random  $k$ -element subset of the vertices of  $G$  contains an independent set of size at least  $ck$  with probability at least  $1 - e^{-ck}$ , where  $c > 0$  is a constant.

- ps6\* H7. Prove that there exists a constant  $c > 0$  so that the following holds. Let  $G$  be a  $d$ -regular graph and  $v_0 \in V(G)$ . Let  $m \in \mathbb{N}$  and consider a simple random walk  $v_0, v_1, \dots, v_m$  where each  $v_{i+1}$  is a uniform random neighbor of  $v_i$ . For each  $v \in V(G)$ , let  $X_v$  be the number times that  $v$  appears among  $v_0, \dots, v_m$ . For that for every  $v \in V(G)$  and  $\lambda > 0$

$$\mathbb{P}\left(\left|X_v - \frac{1}{d} \sum_{w \in N(v)} X_w\right| \geq \lambda + 1\right) \leq 2e^{-c\lambda^2/m}$$

Here  $N(v)$  is the neighborhood of  $v$ .

- H8. Prove that for every  $k$  there exists a  $2^{(1+o(1))k/2}$ -vertex graph that contains every  $k$ -vertex graph as an induced subgraph.

- ps6\* H9. *Tighter concentration of chromatic number*

- (a) Prove that with probability  $1 - o(1)$ , every vertex subset of  $G(n, 1/2)$  with at least  $n^{1/3}$  vertices contains an independent set of size at least  $c \log n$ , where  $c > 0$  is some constant.  
 (b) Prove that there exists some function  $f(n)$  and constant  $C$  such that for all  $n \geq 2$ ,

$$\mathbb{P}(f(n) \leq \chi(G(n, 1/2)) \leq f(n) + C\sqrt{n}/\log n) \geq 0.99.$$

- ps6 H10. Show that for every  $\epsilon > 0$  there exists  $C > 0$  so that every  $S \subseteq [4]^n$  with  $|S| \geq \epsilon 4^n$  contains four elements with pairwise Hamming distance at least  $n - C\sqrt{n}$  apart.

- ps6 H11. *Concentration of measure in the symmetric group.* Let  $U \subseteq S_n$  be a set of at least  $n!/2$  permutations of  $[n]$ . Let  $U_t$  denote the set of permutations that can be obtained starting from some element of  $U$  and then applying at most  $t$  transpositions. Prove that

$$|U_t| \geq (1 - e^{-ct^2/n})n!$$

for every  $t \geq 0$ , where  $c > 0$  is some constant.

Hint in white:

*For the remaining exercises in this section, use Talagrand's inequality*

- H12. Let  $Q$  be a subset of the unit sphere in  $\mathbb{R}^n$ . Let  $\mathbf{x} \in [-1, 1]^n$  be a random vector with independent random coordinates. Let  $X = \sup_{\mathbf{q} \in Q} \langle \mathbf{x}, \mathbf{q} \rangle$ . Let  $t > 0$ . Prove that

$$\mathbb{P}(|X - \mathbb{E}X| \geq t) \leq 4e^{-ct^2}$$

where  $c > 0$  is some constant.

- ps6 H13. *First passage percolation.* Prove that there are constants  $c, C > 0$  so that the following holds. Let  $G$  be a graph, and let  $u$  and  $w$  be two distinct vertices with distance at most  $\ell$  between them. Every edge of  $G$  is independently assigned some random weight in  $[0, 1]$  (not necessarily uniform or identically distributed). The weight of a path is defined to be the sum of the weights of its edges. Let  $X$  be the minimum weight of a path from  $u$  to  $w$  using at most  $\ell$  edges. Prove that there is some  $m \in \mathbb{R}$  so that

$$\mathbb{P}(|X - m| \geq t) \leq Ce^{-ct^2/\ell}.$$

- ps6\* H14. *Second largest eigenvalue of a random matrix.* Let  $A$  be an  $n \times n$  random symmetric matrix whose entries on and above the diagonal are independent and in  $[-1, 1]$ . Show that the second largest eigenvalue  $\lambda_2(A)$  satisfies

$$\mathbb{P}(|\lambda_2(A) - \mathbb{E}\lambda_2(A)| \geq t) \leq Ce^{-ct^2},$$

for every  $t \geq 0$ , where  $C, c > 0$  are constants.

Hint in white:

- H15. *Longest common subsequence.* Let  $(a_1, \dots, a_n)$  and  $(b_1, \dots, b_m)$  be two random sequences with independent entries (not necessarily identically distributed). Let  $X$  denote the length of the longest common subsequence, i.e., the largest  $k$  such that there exist  $i_1 < \dots < i_k$  and  $j_1 < \dots < j_k$  with  $x_{i_1} = y_{j_1}, \dots, x_{i_k} = y_{j_k}$ . Show that, for all  $t \geq 0$ ,

$$\mathbb{P}(X \geq \mathbb{M}X + t) \leq 2 \exp\left(\frac{-ct^2}{\mathbb{M}X + t}\right) \quad \text{and} \quad \mathbb{P}(X \leq \mathbb{M}X - t) \leq 2 \exp\left(\frac{-ct^2}{\mathbb{M}X}\right)$$

where  $c > 0$  is some constant.

## I. ENTROPY METHOD

*The problems in this section should be solved using entropy arguments or results derived from entropy arguments.*

- I1. *Submodularity.* Prove that  $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$ .  
 I2. Let  $\mathcal{F}$  be a collection of subsets of  $[n]$ . Let  $p_i$  denote the fraction of  $\mathcal{F}$  that contains  $i$ . Prove that

$$|\mathcal{F}| \leq \prod_{i=1}^n p_i^{-p_i} (1 - p_i)^{-(1-p_i)}.$$

- ps6\* I3. *Uniquely decodable codes.* Let  $[r]^*$  denote the set of all finite strings of elements in  $[r]$ . Let  $A$  be a finite subset of  $[r]^*$  and suppose no two distinct concatenations of sequences in  $A$  can produce the same string. Let  $|a|$  denote the length of  $a \in A$ . Prove that

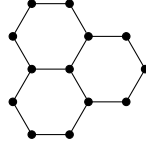
$$\sum_{a \in A} r^{-|a|} \leq 1.$$

- ps6 I4. *Sudoku.* A  $n^2 \times n^2$  *Sudoku square* (the usual Sudoku corresponds to  $n = 3$ ) is an  $n^2 \times n^2$  array with entries from  $[n^2]$  so that each row, each column, and, after partitioning the square into  $n \times n$  blocks, each of these  $n^2$  blocks consist of a permutation of  $[n^2]$ . Prove that the

number of  $n^2 \times n^2$  Sudoku squares is at most

$$\left(\frac{n^2}{e^3 + o(1)}\right)^{n^4}.$$

- ps6 I5. Prove Sidorenko's conjecture for the following graph.



- ps6\* I6. *Triangles versus vees in a directed graph.* Let  $V$  be a finite set,  $E \subseteq V \times V$ , and

$$\Delta = |\{(x, y, z) \in V^3 : (x, y), (y, z), (z, x) \in E\}|$$

(i.e., cyclic triangles; note the direction of edges) and

$$\Lambda = |\{(x, y, z) \in V^3 : (x, y), (x, z) \in E\}|.$$

Prove that  $\Delta \leq \Lambda$ .

- ps6\* I7. *Box theorem.* Prove that for every compact set  $A \subseteq \mathbb{R}^d$ , there exists an axis-aligned box  $B \subseteq \mathbb{R}^d$  with

$$\text{vol } A = \text{vol } B \quad \text{and} \quad \text{vol } \pi_I(A) \geq \text{vol } \pi_I(B) \quad \text{for all } I \subseteq [n].$$

Here  $\pi_I$  denotes the orthogonal projection onto the  $I$ -coordinate subspace.

(For the purpose of the homework, you only need to establish the case when  $A$  is a union of grid cubes. It is optional to give the limiting argument for compact  $A$ .)

- I8. Let  $\mathcal{G}$  be a family of graphs on vertices labeled by  $[2n]$  such that the intersection of every pair of graphs in  $\mathcal{G}$  contains a perfect matching. Prove that  $|\mathcal{G}| \leq 2^{\binom{2n}{2} - n}$ .
- I9. *Loomis–Whitney for sumsets.* Let  $A, B, C$  be finite subsets of some abelian group. Writing  $A + B := \{a + b : a \in A, b \in B\}$ , etc., prove that

$$|A + B + C|^2 \leq |A + B| |A + C| |B + C|.$$

- ps6\* I10. *Shearer for sums.* Let  $X, Y, Z$  be independent random integers. Prove that

$$2H(X + Y + Z) \leq H(X + Y) + H(X + Z) + H(Y + Z).$$

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18.226 Probabilistic Methods in Combinatorics  
Fall 2022

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